# Behavior-based Access Control in Wired and Wireless Networks

Angelos D. Keromytis Network Security Lab Department of Computer Science

## Network Access Control (NAC)

- Mechanism used to control access to a network
- Complements a firewall
  - firewalls keep untrusted nodes out
  - NAC allows trusted nodes to connect

#### Motivation: Pre-connect Phase



#### Motivation: Pre-connect Phase



#### Motivation: Pre-connect Phase



#### Motivation: Post-connect Phase



Computer Science at Columbia University

#### Motivation: Post-connect Phase



Computer Science at Columbia University



Computer Science at Columbia University







Computer Science at Columbia University

## Related Work

- Admission and Access control in NAC technologies
  - PacketFence, Cisco NAC, Microsoft NAP, Symantec
  - All pre- and post-connect policies are manually determined
- Admission and Access control in MANETS
  - Threshold cryptography: no specification of how decision is made (Narasimha et al. '03,Yin et al. '07)
  - IDS used for access control at a routing level (AODV, DSR) not application level (Parker et al. '04, Huang et al. '03)

### Problem Statement

• How to automate the creation and update of admission and access control policies in a network?

• We seek a solution that should be applicable to both centralized network architectures (NAC) and distributed network architectures like MANETs.

#### Solution: Behavior-based Admission and Access Control

- Derive admission and access control policies by profiling the behavior of network devices.
- Behavior profiles modeled by an AD sensor will be used to automatically identify what constitutes normal behavior within the network.
- Behavior is defined as the payload or volumetric characteristics of the network traffic exchanged.

#### Solution: Behavior-based Admission and Access Control

- Derive admission and access control policies by profiling the behavior of network devices.
- Behavior profiles modeled by an AD sensor will be used to automatically identify what constitutes normal behavior within the network.
- Behavior is defined as the payload or volumetric characteristics of the network traffic exchanged.

De Facto vs. De Jure policies

### Contributions

• Enhancement by automating the creation of behavior-based policies for admission and access control

- Automatic and robust update of behavior-based policies as profiles evolve over time
- An approach that is applicable to both centralized and fully distributed networks

### Outline

- Behavior-based NAC Mechanism
- Automatic Clustering and Policy Update
- Cluster-based AD sensor
- Behavior-based Mechanism for MANETs

• Each device has an AD sensor to model,

Input behavior profile

Output behavior profile

Bad profile: malware knowledge

All are automatically updated by the AD sensor

- Each device has an AD sensor to model,
  - Input behavior profile
  - Output behavior profile
  - Bad profile: malware knowledge
  - All are automatically updated by the AD sensor
- BB-NAC Phases:
  - Initial Setup
  - Admission Control (Pre-connect Phase)
  - Access Control (Post-connect Phase)

• Each device has an AD sensor to model,

Input behavior profile

Output behavior profile

Bad profile: malware knowledge

All are automatically updated by the AD sensor

• BB-NAC Phases:

Initial Setup

Admission Control (Pre-connect Phase)

Access Control (Post-connect Phase)

 NAC enforcer and high-ranked monitors responsible for the admission and access control

### **BB-NAC: Initial Setup**



Computer Science at Columbia University



#### Network Security Lab

### Initial Setup

- Pre-determined clusters of behavior profiles:
  - Clients and servers (per port and direction)
  - Each cluster of behavior profiles represents a valid

behavior for the network

• For each cluster member, NAC enforcer determines a threshold or local boundary of normal behavior,

 $t_{P_i} = \max_{j=0..n} d(P_i, P_j)$  for each profile  $P_i$  in cluster

• Each device is represented by  $M_i = \{P_i, t_{P_i}, B_i\}$ 

Network Security Lab

### Admission Control





Computer Science at Columbia University

• New device with behavior profile  $P_{new}$  and bad profile  $B_{new}$ 

- New device with behavior profile  $P_{new}$  and bad profile  $B_{new}$  solicits admission to network

- New device with behavior profile  $P_{new}$  and bad profile  $B_{new}$  solicits admission to network
- Behavior Profile Check

► Is the behavior of the device similar to any cluster of behavior in the network?

- New device with behavior profile  $P_{new}$  and bad profile  $B_{new}$  solicits admission to network
- Behavior Profile Check

► Is the behavior of the device similar to any cluster of behavior in the network?

- Bad Profile Check
  - Does the device have enough malware knowledge?

- New device with behavior profile  $P_{new}$  and bad profile  $B_{new}$  solicits admission to network
- Behavior Profile Check

► Is the behavior of the device similar to any cluster of behavior in the network?

• Bad Profile Check

Does the device have enough malware knowledge?

• If admission rejected multiple times, device IP is blacklisted

Computer Science at Columbia University

• Group decision based on the sum of individual decisions from each member of the closest cluster (percentage of agreement)

 $cluster = \min_{i=0..n} d(c[i], P_{new})$ 

- Group decision based on the sum of individual decisions from each member of the closest cluster (percentage of agreement)
- Individual decisions evaluate the difference between their local behavior and the newcomer's behavior

 $cluster = \min_{i=0..n} d(c[i], P_{new})$ 

 $v_i = 1$  if  $d(P_i, P_{new}) \le t_{P_i}$  where  $P_i \in cluster$ 

- Group decision based on the sum of individual decisions from each member of the closest cluster (percentage of agreement)
- Individual decisions evaluate the difference between their local behavior and the newcomer's behavior

 $cluster = \min_{i=0..n} d(c[i], P_{new})$ 

 $v_i = 1$  if  $d(P_i, P_{new}) \le t_{P_i}$  where  $P_i \in cluster$ 

- Group decision based on the sum of individual decisions from each member of the closest cluster (percentage of agreement)
- Individual decisions evaluate the difference between their local behavior and the newcomer's behavior

$$cluster = \min_{i=0..n} d(c[i], P_{new})$$

 $v_i = 1$  if  $d(P_i, P_{new}) \le t_{P_i}$  where  $P_i \in cluster$ 

- Group decision based on the sum of individual decisions from each member of the closest cluster (percentage of agreement)
- Individual decisions evaluate the difference between their local behavior and the newcomer's behavior

Closest Cluster: 
$$cluster = \min_{i=0..n} d(c[i], P_{new})$$
Individual decision:  $v_i = 1$  if  $d(P_i, P_{new}) \le t_{P_i}$  where  $P_i \in cluster$ 

Computer Science at Columbia University
# Behavior Profile Check

- Group decision based on the sum of individual decisions from each member of the closest cluster (percentage of agreement)
- Individual decisions evaluate the difference between their local behavior and the newcomer's behavior

$$cluster = \min_{i=0..n} d(c[i], P_{new})$$

- Individual decision:
- Final Group Decision:

$$v_i = 1$$
 if  $d(P_i, P_{new}) \le t_{P_i}$  where  $P_i \in cluster$   
 $v = \frac{1}{n} \sum_{i=0..n} v_i$  if  $v \succ \varepsilon$  ACCEPT

#### Pre-connect Phase: Example



Computer Science at Columbia University

#### Pre-connect Phase: Example



#### Pre-connect Phase: Example











## Bad Profile Check

Computer Science at Columbia University

# Bad Profile Check

• Group decision: each member evaluates malware knowledge of newcomer

$$v_i = 1$$
 if  $B_i \subset B_{new}$  where  $B_i \in cluster$ 

# Bad Profile Check

• Group decision: each member evaluates malware knowledge of newcomer

Individual decision:
 
$$v_i = 1$$
 if  $B_i \subset B_{new}$  where  $B_i \in cluster$ 
 Final Group Decision:
  $v = \frac{1}{n} \sum_{i=0..n} v_i$  if  $v \succ \varepsilon$  ACCEPT

## Access Control

- Traffic is deemed normal or anomalous based on a group-profile decision
- All members in a cluster vote, not only the ones exchanging traffic: group knowledge

$$v = \frac{1}{n} \sum_{k=1}^{n} P_{k,d}(g)$$
 if  $v \succ \varepsilon$  ANOMALY

- If anomaly detected, generate alert / quarantine device
- Anomalous traffic is used to update bad profiles
- Normal traffic is used to train behavior profiles

#### Proof-of-concept Experiments

- Four CS webservers (servers cluster, port 80)
- ANAGRAM AD sensor: input content profiles (BFs)
- Initial Setup
  - Stable profiles computed with two weeks of traffic
- Admission Control (FR/TR)
  - Normal Testing Set
    - Stable profiles computed for the following two weeks
  - Malicious Testing Set
    - Profiles for 8 DNS servers computed for same two weeks
    - Normal profiles poisoned with one virus from vxheavens
- Access Control (FP/DR)
  - Traffic from worms: CodeRed, Webdav, Mirela, php worms

20

Scenario	FR(%)	$TR_t$ (%)	$TR_v$ (%)	
(i) 25%	0	100	0	
(ii) 50%	0	100	0	
(iii) 75%	0	100	0	
(iv) 100%	0	100	0	

#### **Profiles port 53**

Scenario	FR(%)	$TR_t$ (%)	$TR_v$ (%)			
(i) 25%	0	100	0			
(ii) 50%	0	100	0			
(iii) 75%	0	100	0			
(iv) 100%	0	100	0			

#### **Profiles poisoned**

Scenario	FR(%)	$TR_t$ (%)	$TR_v$ (%
(i) 25%	0	100	0
(ii) 50%	0	100	0
(iii) 75%	0	100	0
(iv) 100%	0	100	0

Admission Control

Scenario	FR(%)	$TR_t$ (%)	$TR_v$ (%)	
(i) 25%	0	100	0	
(ii) 50%	0	100	0	
(iii) 75%	0	100	0	
(iv) 100%	0	100	0	

Access Control

(iv) 100%	83%	0.001%	server4	99%	0.01%
(iii) 75%	99%	0.005%	server3	99%	0.015%
(ii) 50%	99%	0.02%	server2	83%	0.009%
(i) 25%	100%	0.032%	server1	100%	0.02%
Percentage	DR	FP	Server	DR	FP

Science at

**J**niversity

Admission Control

Scenario	FR(%)	$TR_t$ (%)	$TR_v$ (%)	
(i) 25%	0	100	0	
(ii) 50%	0	100	0	
(iii) 75%	0	100	0	
(iv) 100%	0	100	0	

Access Control

#### **Best Collaborative Solution** *E* **=**75%

	21							(	Columl	bia Unive
	Group Rates			Ine	divic	lual	Rat	es	ter Scier	
	(iv) 100%	83%	0.001%			server4	99%	0.01%		
<	(iii) 75%	99%	0.005%	>		server3	99%	0.015%		
	(ii) 50%	99%	0.02%			server2	83%	0.009%		
	(i) 25%	100%	0.032%			server1	100%	0.02%		
	Percentage	DR	$\mathbf{FP}$			Server	DR	$\mathbf{FP}$		
					22					

ce at

rsity

#### Concept Drift in Behavior Profiles

#### Concept Drift in Behavior Profiles

 $P_i = P_i \wedge P_{i-1} \wedge \ldots \wedge P_{i-q}$ 

#### Concept Drift in Behavior Profiles

 $P_i = P_i \wedge P_{i-1} \wedge \ldots \wedge P_{i-q}$ 

#### Concept Drift in Behavior Profiles

• AllModels: keep all previous knowledge

$$P_i = P_i \land P_{i-1} \land \dots \land P_{i-q}$$

#### Concept Drift in Behavior Profiles

• AllModels: keep all previous knowledge

$$P_i = P_i \land P_{i-1} \land \dots \land P_{i-q}$$

#### Concept Drift in Behavior Profiles

- AllModels: keep all previous knowledge
- ANDModels: keep only common knowledge between models

$$P_i = P_i \wedge P_{i-1} \wedge \dots \wedge P_{i-q}$$

#### Concept Drift in Behavior Profiles

- AllModels: keep all previous knowledge
- ANDModels: keep only common knowledge between models

$$P_i = P_i \land P_{i-1} \land \dots \land P_{i-q}$$

• Add new behaviors, OR ANDModels:

$$P_i = \bigvee_{t=0}^{s-1} P_{i-t} \wedge P_{i-1-t} \wedge \dots \wedge P_{i-q-t}$$

#### Concept Drift in Behavior Profiles



Computer Science at Columbia University

#### Concept Drift in Behavior Profiles



Computer Science at Columbia University

• Pre-connect latency

$$l = l_a + (1 - \rho) \times l_q$$

• Pre-connect latency

$$l = l_a + (1 - \rho) \times l_q$$

 $l = (1 - FP) \times l_{BF} + FP \times l_q$ 

• Pre-connect latency

$$l = l_a + (1 - \rho) \times l_q$$

Post-connect latency

$$l = (1 - FP) \times l_{BF} + FP \times l_q$$

• Pre-connect latency

$$l = l_a + (1 - \rho) \times l_q$$

Post-connect latency

$$l = (1 - FP) \times l_{BF} + FP \times l_q$$

• Pre-connect latency

$$l = l_a + (1 - \rho) \times l_q$$

Post-connect latency

$$l = (1 - FP) \times l_{BF} + FP \times l_q$$

BB-NAC values (cluster of 10 devices):

 Pre-connect: 180~342ms
 Post-connect (Best collaborative): 5,785~50.56ms

## Summary of Results

New mechanism to automatically create *De Facto* admission and access control policies

A novel admission and access control based on a profile-group decision process which may outperform individual decision processes

Proof-of-concept evaluation using content behavior profiles (hashed into BFs) from 4 webservers

## Outline

- Behavior-based NAC technologies
   —Automatic Clustering and Policy Update
- Cluster-based AD sensor
- Behavior-based Policies for MANETs
- Conclusions and Future Work

# Automatic Clustering and Policy Update \*

- Automatic computation of clusters of behavior

   Initial setup complemented with *k-means clustering* Other phases proceed normally
- Robust evolution of clusters of behavior over time
   Incremental-learning algorithm: differentiate between concept drift and attack
- NAC enforcer is responsible for both tasks



### CLUSTERING

- Initially all devices in network communicate behavior profile to NAC enforcer
- K-means clustering is performed on a per port, per direction basis to identify clusters of similar behavior
- Although computationally expensive, is performed only once during the setup
- Best number of clusters is selected via cross-validation based on best performance for the access control

## CLUSTERING

• Behavior Profiles are of the type

$$p_i = \{p_i[0], p_i[1], \dots p_i[n]\}$$

• Each field  $p_i[\ell]$  represents the average measure of any volumetric characteristic of the connections established by a user

• Distances between profiles: euclidean distance

$$d(p_i, p_j) = \sqrt{(\sum_{\ell=0..n} (p_i[\ell] - p_j[\ell])^2)}$$

• Normalization of measurements to avoid distance calculation dominated by measurements at different scales
# Clustering

- Kmeans++: smart selection of initial seeds (non-deterministic)
- Cross-validation determines best values to be used during admission and access control, r = (1-FR) + TR

Best k distribution

 $\blacktriangleright$  Best percentage of agreement  $\mathcal{E}$ 

$$v = \frac{1}{n} \sum_{i=0..n} v_i$$

Veighted or non-weighted voting (w)  

$$W_i = \frac{d_{\text{max}} - d_i}{d_{\text{max}} - d_{\text{min}}} \times 1$$

Computer Science at Columbia University

# BOOTSTRAP

- NAC enforcer computes distance between each profile in the cluster and the others
- This measurement represents a threshold used during access control to determine which devices can be accepted into the network

# ACCESS CONTROL

•Upon arrival, newcomer presents its profile

•NAC enforcer computes closest behavioral cluster

•NAC enforcer computes a "voting process" across profiles in closest cluster

• Members vote for or against with thresholds computed during bootstrap phase

- Simple voting: all votes count equally. If 50% or more agree on normalcy, device is accepted
- Weighted voting: profiles closer to newcomer have a stronger vote

### EXAMPLE



# Policy Update

- Over time, new behavior profiles may be computed
  - Admitted as normal by one of the existing clusters
  - New behavior evolved from old behavior (concept drift)
  - Malicious behavior profile (attack)

 Behaviors are defined as clusters
 New behavior accepted if defines a cluster within the boundaries of existing behaviors

### Incremental-Learning Algorithm

### decision boundary



### INCREMENTAL LEARNING ALGORITHM 2



## Attacks against the Mechanism

### Collusion Attacks

- Attackers try to craft malicious behavior profiles to create a new cluster
- Threshold Attacks
  - Attackers try to modify thresholds to change dimensions of own cluster
  - Unified: stretch in one direction
  - Diversified: stretch in multiple directions
- Experimental results show robustness against both types of attacks

Ground Truth

Computer Science at Columbia University

- Cisco NetFlow logs from router at Columbia University
- 300 Columbia IPs (hosts) were randomly selected for port 80
  - 300 Profiles computed with their flows (week1)
  - 300 Profiles computed with their flows (week2)
  - Each behavior characterized 7 flow parameters:
    - total number of flows
    - average flow size

- Cisco NetFlow logs from router at Columbia University
- 300 Columbia IPs (hosts) were randomly selected for port 80
  - 300 Profiles computed with their flows (week1)
  - 300 Profiles computed with their flows (week2)
  - Each behavior characterized 7 flow parameters:
    - total number of flows
    - average flow size
    - average flow duration

- Cisco NetFlow logs from router at Columbia University
- 300 Columbia IPs (hosts) were randomly selected for port 80
  - 300 Profiles computed with their flows (week1)
  - 300 Profiles computed with their flows (week2)
  - Each behavior characterized 7 flow parameters:
    - total number of flows
    - average flow size
    - average flow duration
    - total number of packets contained in all flows

Computer Science at Columbia University

**Ground Truth** 

- Cisco NetFlow logs from router at Columbia University
- 300 Columbia IPs (hosts) were randomly selected for port 80
  - 300 Profiles computed with their flows (week1)
  - 300 Profiles computed with their flows (week2)
  - Each behavior characterized 7 flow parameters:
    - total number of flows
    - average flow size
    - average flow duration
    - total number of packets contained in all flows
    - average number of packets per flow

- Cisco NetFlow logs from router at Columbia University
- 300 Columbia IPs (hosts) were randomly selected for port 80
  - 300 Profiles computed with their flows (week1)
  - 300 Profiles computed with their flows (week2)
  - Each behavior characterized 7 flow parameters:
    - total number of flows
    - average flow size
    - average flow duration
    - total number of packets contained in all flows
    - average number of packets per flow
    - total number of unique IP addresses contained in all flows

- Cisco NetFlow logs from router at Columbia University
- 300 Columbia IPs (hosts) were randomly selected for port 80
  - 300 Profiles computed with their flows (week1)
  - 300 Profiles computed with their flows (week2)
  - Each behavior characterized 7 flow parameters:
    - total number of flows
    - average flow size
    - average flow duration
    - total number of packets contained in all flows
    - average number of packets per flow
    - total number of unique IP addresses contained in all flows
    - average packet size

Computer Science at Columbia University



- average flow duration
- total number of packets contained in all flows
- average number of packets per flow
- total number of unique IP addresses contained in all flows
- average packet size

Ground Truth

- total number of packets contained in all flows
- average number of packets per flow
- total number of unique IP addresses contained in all flows
- average packet size

Ground Truth

- average number of packets per flow
- total number of unique IP addresses contained in all flows
- average packet size



### • total number of unique IP addresses contained in all flows

average packet size

Computer Science at Columbia University

Ground Truth

average packet size

Computer Science at Columbia University

- Simulated NAC with 300 members, other 300 attempt admission
- Training Set: 300 profiles week1
- Cross-validation Set: 75 profiles week2 (FR)
- ► Testing Set: 225 profiles week2 (FR)
- ► TR: profiles at one, two and three  $\sigma$  away from individual clusters

### VALIDATION EXPERIMENTS: PHASES

- Simulated NAC with 300 members, other 300 attempt admission
- Training Set: 300 profiles week1
- · Cross-validation Set: 75 profiles week2 (FR)
- Testing Set: 225 profiles week2 (FR)
- TR: create artificial profiles at one, two and three standard deviations away from individual clusters

### VALIDATION EXPERIMENTS: CLUSTERING AND CROSS-VALIDATION



Best Performance: K=40 clusters (TR=95%, FP=10%)

Computer Science at Columbia University

### VALIDATION EXPERIMENTS: ACCESS CONTROL

- @one and two standard deviation some anomalous profiles go undetected due to problems with the profile generation itself
- Future work will evaluate real anomalous profiles and not artificially generated

$\sigma$ From Individual Clusters	True Rejection Rate
$1 \sigma$	95%
$2 \sigma$	98%
$3\sigma$	100%

### VALIDATION EXPERIMENTS: CONCEPT DRIFT AND COLLUSION ATTACKS

- Create anomalous profiles at one, two and three standard deviations away from the global behavioral centroid
- Idea: create outlier profiles to determine the boundary the algorithm creates between concept drift and attacks
- Important: the creation of new clusters is limited by the algorithm so that the damage an attacker can infer is also limited

$\sigma$ From Global Centroid	Candidate Clusters Rejected
$1 \sigma$	85%
$2 \sigma$	92%
$3\sigma$	96%

ence at

### VALIDATION EXPERIMENTS: THRESHOLD ATTACKS

- Selected two clusters from clustering phase: lowest and highest spread (distance across behaviors)
- Attacks: each attacker generates a different profile and iteratively increases feature values by 10%
  - Diversified attack: each attacker generates a different profile
  - Unified attack: all initial profiles are equal

• Dispersion: ratio between initial average threshold and final average threshold after attack

### VALIDATION EXPERIMENTS: THRESHOLD ATTACKS

- All attacks are eventually detected by the mechanism
- Dispersion factor increased at maximum by three times



## Summary of Results



Mechanism to automatically determine clusters of behavior and other parameters specific to the network



Mechanism is robust over time while allowing for creation of new behaviors (concept drift)



Mechanism can detect collusion attacks and limits the damage incurred by threshold attacks



Extensive evaluation using volumetric profiles (flows) of 300 hosts at Columbia University

### Outline

- Behavior-based NAC technologies
- Automatic Clustering and Policy Update
  Cluster-based AD sensor
- Behavior-based Policies for MANETs
- Conclusions and Future Work

### Sensor Details

### Single-Profile AD sensor

#### **Cluster-based AD sensor**



### Sensor Details

### Single-Profile AD sensor

#### **Cluster-based AD sensor**



### Sensor Details

### Single-Profile AD sensor

### **Cluster-based AD sensor**



$$D_{P_x}(t) = \begin{cases} 0, & \text{if } P_x - \sigma \prec t \prec P_x + \sigma \\ 1, & \text{otherwise} \end{cases}$$

$$S_{x} = \begin{cases} \sum_{i=1}^{x, y, z} D_{P_{i}}(t) \\ 1, \text{ if } \frac{\sum_{i=1}^{x, y, z} D_{P_{i}}(t)}{3} \ge \varepsilon \\ 0, \text{ otherwise} \end{cases}$$

- CRAWDAD repository: 1 month of tcpdump wireless traffic (no content)
- Ports: 21 (FTP), 22 (SSH), 25 (SMTP), 80 (HTTP)
- For each port, we identified 100 different MAC addresses with output traffic to the services
- Behavior profiles for each user/service were computed as daily histograms of usage for the first week of data

$$P_i = \{h_{f_1}, \dots, h_{f_n}\} \text{ where } h_{f_n} = \{(\alpha_0, \sigma_0), (\alpha_1, \sigma_1), \dots, (\alpha_{23}, \sigma_{23})\}$$

- Parameters modeled:
  - average number of unique users contacted per hour (i)
  - average number of packets exchanged per hour (ii)
  - average length of the packets exchanged per hour (iii)
- Performance measured with average false positive  $\overline{FP}$  and average detection threshold  $\bar{\phi}$



# Summary of Results

Cluster-based AD sensor provides a broader definition of normal behavior that compensates for poor or insufficient training of individual sensors and reduces the volume of false alerts

### Outline

- Behavior-based NAC technologies
- Automatic Clustering and Policy Update
- Cluster-based AD sensor
- Behavior-based Policies for MANETs
- Conclusions and Future Work
Does the BB-NAC mechanism work for MANETs?

Does the BB-NAC mechanism work for MANETs?
 Automatically derive policies from behavior profiles

- Does the BB-NAC mechanism work for MANETs?
  - Automatically derive policies from behavior profiles
  - Fully distributed approach to profile computation and alert generation/analysis

- Does the BB-NAC mechanism work for MANETs?
  Automatically derive policies from behavior profiles
  Fully distributed approach to profile computation and alert generation/analysis
- BARTER: adaptation of BB-NAC to fully distributed environments

- Does the BB-NAC mechanism work for MANETs?
  Automatically derive policies from behavior profiles
  - Fully distributed approach to profile computation and alert generation/analysis
- BARTER: adaptation of BB-NAC to fully distributed environments
  - Threshold cryptographic layer (t,n) for fully distributed management

- Does the BB-NAC mechanism work for MANETs?
  Automatically derive policies from behavior profiles
  Fully distributed approach to profile computation and aler
  - Fully distributed approach to profile computation and alert generation/analysis
- BARTER: adaptation of BB-NAC to fully distributed environments
  - Threshold cryptographic layer (t,n) for fully distributed management
  - Light-weight version. Instead of clusters of behavior, MANET devices derive thresholds from their top t-1 most similar profiles

- Does the BB-NAC mechanism work for MANETs?
  Automatically derive policies from behavior profiles
  - Fully distributed approach to profile computation and alert generation/analysis
- BARTER: adaptation of BB-NAC to fully distributed environments
  - Threshold cryptographic layer (t,n) for fully distributed management
  - Light-weight version. Instead of clusters of behavior, MANET devices derive thresholds from their top t-1 most similar profiles
  - Members only exchange output behavior profiles to avoid

- Does the BB-NAC mechanism work for MANETs?
  Automatically derive policies from behavior profiles
  Fully distributed approach to profile computation and alert
  - generation/analysis
- BARTER: adaptation of BB-NAC to fully distributed environments
  - Threshold cryptographic layer (t,n) for fully distributed management
  - Light-weight version. Instead of clusters of behavior, MANET devices derive thresholds from their top t-1 most similar profiles
  - Members only exchange output behavior profiles to avoid crafted attacks adapted to input profiles

#### Network Security Lab

## BARTER SETUP

• Devices initially exchange their output behavior profiles and compute their local thresholds:

 $t_{P_i} = \max_{j=0..n} d(P_{i,in}, P_{j,out})$  where  $P_j \in \text{top t-1 most similar}$ 

• Each member computes its own threshold as,



• Cross-Validation to determine best t/n for the network

## BARTER CRYPTO SETUP

- Shamir Secret Sharing to divide the secret into n shares
- Each member computes its partial shared key such that any *t* members will be able to recover the secret and generate new keys for new members

## BARTER ADMISSION CONTROL

• Individual admission control decision:

 $v_i = 1$  if  $d(P_{i,in}, P_{new,out}) \le \tau_i$ 

- Group decision: if *t* MANET members agree, newcomer accepted
- Newcomer receives output profiles from all other members and creates its local table
- To avoid recalculating t/n upon every new accepted member:

• t is updated only when the ratio exceeds the range (t/n+w, t/n-w)

## BARTER CRYPTO ADMISSION

- Newcomer broadcasts its public key certificate and its behavior profile
- Each individual member that accepts the newcomer as normal, generates a partial signature.
- Newcomer computes signature by summing *t* partial signatures received upon acceptance.

## BARTER ACCESS CONTROL

- Each member AD is continuously screening traffic
- If an anomaly is detected, it is shared with its *top t-1* most similar members
- If the *t-1* members agree on the anomaly, the device is expelled
- Assumption: existence of scheme to avoid data tampering and prevent devices from falsifying alerts or replay attacks

## BARTER CRYPTO ACCESS

- Proactive Key Generation is used to eliminate a device from accessing further communications
- The "bad" device is added to a Certificate Revocation List (CRL) which is broadcasted to all members
- Devices that receive t CRLs renew keys through point-topoint encrypted communication channels with members not in the CRL

### BARTER ATTACKS

- DDoS costs: robustness against DDoS attacks
- Normalized value (0..1)

$$DDoS \begin{cases} 0.5, \text{ if } \frac{t}{n} = 0.5\\ \frac{t}{n}, \text{ if } \frac{t}{n} \prec 0.5\\ 1 - \frac{t}{n}, \text{ if } \frac{t}{n} \succ 0.5\\ 1 - \frac{t}{n}, \text{ if } \frac{t}{n} \succ 0.5 \end{cases}$$

Computer Science at Columbia University

## BARTER CRYPTOGRAPHIC COSTS

• Cryptographic costs incurred in key (re)generation upon admission and access control (when t is (in/de)cremented)

$$CC = K \times \sum_{n_0}^{n_{final}} (update \times n) - 1$$
$$update = \begin{cases} 1, & \text{if } \frac{t}{n} \prec \left(\frac{t_0}{n_0} - w\right) & \text{or } \text{if } \frac{t}{n} \succ \left(\frac{t_0}{n_0} + w\right) \\ 0, & \text{otherwise} \end{cases}$$

- 140 users from ENRON dataset, email-like MANET application
- Computed input and output profiles using Shanner's Algorithm

$$W(i) = \log \frac{x_i}{N_g} \times \frac{1}{\log N_g} \sum_{j=1}^{N_g} p_{ij} \log \frac{1}{p_{ij}} \times \left(1 - \frac{1}{\log L}\right)^{goodS, badS} \sum_{j=1}^{goodS, badS} p_{ij} \log \frac{1}{p_{ij}}$$

- Good samples: ENRON emails
- Bad samples: Signature content of Snort rules (58) and 600 virus samples from vxheavens

- 140 users from ENRON dataset, email-like MANET application
- Computed input and output profiles using Shanner's Algorithm

$$W(i) = \log \frac{x_i}{N_g} * \frac{1}{\log N_g} \sum_{j=1}^{N_g} p_{ij} \log \frac{1}{p_{ij}} \times \left(1 - \frac{1}{\log L}\right)^{goodS, badS} \sum_{j=1}^{goodS, badS} p_{ij} \log \frac{1}{p_{ij}}$$

- Good samples: ENRON emails
- Bad samples: Signature content of Snort rules (58) and 600 virus samples from vxheavens

- 140 users from ENRON dataset, email-like MANET application
- Computed input and output profiles using Shanner's Algorithm

$$W(i) = \log \frac{x_i}{N_g} * \frac{1}{\log N_g} \sum_{j=1}^{N_g} p_{ij} \log \frac{1}{p_{ij}} \times \left(1 - \frac{1}{\log L}\right)^{goodS, badS} p_{ij} \log \frac{1}{p_{ij}}$$

- Good samples: ENRON emails
- Bad samples: Signature content of Snort rules (58) and 600 virus samples from vxheavens

- 140 users from ENRON dataset, email-like MANET application
- Computed input and output profiles using Shanner's Algorithm

$$W(i) = \log \frac{x_i}{N_g} * \left[ \frac{1}{\log N_g} \sum_{j=1}^{N_g} p_{ij} \log \frac{1}{p_{ij}} \right] * \left( 1 - \frac{1}{\log L} \frac{1}{j} \sum_{j=1}^{goodS, badS} p_{ij} \log \frac{1}{p_{ij}} \right]$$

- Good samples: ENRON emails
- Bad samples: Signature content of Snort rules (58) and 600 virus samples from vxheavens

- Behavior Profiles are saved as Bloom Filters
- BFs are one-way structures that preserve the privacy of the output content of the users
- BARTER saves its behavior profiles (input and output) as BFs
- Distance between profiles is computed as XOR that quantifies the amount of entries that differ across profiles:

$$d(P_i, P_j) = |P_i \oplus P_j|$$

- Training Set: 80 ENRON behavior profiles  $\rightarrow$  Ground truth
- Cross-validation Set:
  - 30 ENRON behavior profiles (ground truth)
  - 30 code/executables-based behavior profiles
- Testing Set:
  - 30 behavior ENRON profiles (ground truth)
  - 30 code/executables-based behavior profiles
- Admission Control measured in terms of TR and FR rates
- Cross-validation
  - For each t/n we compute r = (1-FR) + TR + (1-CC) + DDoS
  - The highest ranked t/n is selected for admission and access control
  - Experiments run multiple times and results are averaged

## EXPERIMENTAL EVALUATION $\begin{array}{l} & CONTENT \ PROFILES \\ & r = (1 - FR) + TR + (1 - CC) + DD_0S \end{array}$



Ratio $t/n$	Absolute Cost	Normalized Cost	DDoS
$0.1\pm0.02$	175  imes K	0.07	0.1
$0.3\pm0.02$	$934 \times K$	0.37	0.3
$0.5\pm0.02$	$1576 \times K$	0.63	0.5
$0.7\pm0.02$	$2146 \times K$	0.86	0.3
$0.9\pm0.02$	$2484 \times K$	1.0	0.1

Computer Science at Columbia University

## EXPERIMENTAL EVALUATION $\begin{array}{l} & CONTENT \ PROFILES \\ & r = (1 - FR) + TR + (1 - CC) + DDoS \end{array}$



## EXPERIMENTAL EVALUATION $\begin{array}{l} & CONTENT \ PROFILES \\ & r = (1 - FR) + TR + (1 - CC) + DDoS \end{array}$



Best ratio t/n=10% ► Admission Control Results: FR= 0.13 TR= 1

> Computer Science at Columbia University

Monday, June 28, 2010

#### EXPERIMENTAL EVALUATION VOLUMETRIC PROFILES

- 140 users from ENRON dataset
- Computed input and output profiles as daily histograms on the
- number of emails exchanged (EMT tool),

$$P_{i,d} = \{b_1, ..., b_{bg}\}$$
 where  $b_j = \{(\alpha, \sigma)\}$ 

- Two types of volumetric profiles:
  - Hourly Histrograms: bg=24
  - Grouped Histograms: bg=4
  - Profile distance computed using Euclidean distance
  - Grouped Histograms save more bandwidth during its exchange
  - Best profiling technique is selected during cross-validation, together with best t/n % f(x) = 0

#### EXAMPLES OF A VOLUMETRIC PROFILE



Hourly Profile bg=24

Grouped Profile bg=4

Computer Science at Columbia University

#### EXPERIMENTAL EVALUATION VOLUMETRIC PROFILES

- Simulated MANET running an email-like application
- Admission Control quality measured in terms of TR and FR
- Good Samples: ENRON volumetric profiles
- Bad Samples: Synthetically computed as profiles that are one, two and three standard deviations away from the top t-1 entries in the local table of each MANET member.
- Training Set (80), Cross-validation Set (30) and Testing Set (30)
- Cross-validation determines ratio t/n, type of profile

# $\begin{array}{l} EXPERIMENTAL\\ EVALUATION\\ r=(1-FR)+TR+(1-CC)+DDoS \end{array}$



Ratio $t/n$	CC for Hourly Histograms	CC for Grouped Histograms	DDoS
$0.1\pm0.02$	155  imes K	$157 \times K$	0.1
$0.3\pm0.02$	913  imes K	915  imes K	0.3
$0.5\pm0.02$	$1564 \times K$	$1570 \times K$	0.5
$0.7\pm0.02$	$2131 \times K$	$2140 \times K$	0.3
$0.9\pm0.02$	$2581 \times K$	$2590 \times K$	0.1

Computer Science at Columbia University Network Security Lab



Network Security Lab



 Best ratio t/n=50%, and grouped histogram
 ▶ Admission Control Results: FR= 0.08 TR= 0.9

> Computer Science at Columbia University

## Summary of Results



Adaptation of BB-NAC to MANETs



Integration of the mechanism with a threshold cryptographic layer to achieve fully distributed decisions



Evaluation of BARTER with content and volumetric profiles from ENRON dataset

## Outline

- Behavior-based NAC technologies
- Automatic Clustering and Policy Update
- Cluster-based AD sensor
- Behavior-based Policies for MANETs
- Conclusions and Future Work

## **Conclusions / Contributions**



Novel mechanism to automatically create and update behavior-based policies for NAC technologies and MANETs



Incremental-Learning Algorithm that makes the mechanism robust against attacks from network members



Cluster-based AD sensor that reduces volume of false alerts by sharing the detection with similar behavior profiles



Extensive experimental evaluation of BB-NAC and BARTER with different types of behavior profiles and datasets